

# measurement sharpness and incompatibility as quantum resources

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*Probability theory is a measure theory with a soul. (Mark Kac)*

*Quantum foundations is a linear algebra with a soul. (anon.)*

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## references

- F.B., K. Kobayashi, S. Minagawa:  
*A complete and operational resource theory of measurement sharpness.*  
Arxiv:2303.07737
- F.B., K. Kobayashi, S. Minagawa, P. Perinotti, A. Tosini:  
*Unifying different notions of quantum incompatibility into a strict hierarchy of resource theories of communication.*  
Quantum 7, 1035 (2023)

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## POVMs and instruments

in this talk: all sets ( $\mathbb{X}, \mathbb{Y}$  etc.) are finite, all spaces ( $\mathcal{H}_A, \mathcal{H}_B$  etc.) are finite-dimensional

**POVM**: family  $\mathbf{P}$  of positive semidefinite operators on  $\mathcal{H}$  labeled by set  $\mathbb{X}$ , i.e.,  $\mathbf{P} = \{P_x\}_{x \in \mathbb{X}}$ , with  $P_x \geq 0$  and  $\sum_x P_x = \mathbb{1}$

**interpretation**: expected probability of outcome  $x$  is  $p(x) = \text{Tr}[\varrho P_x]$

**instrument**: family  $\{\mathcal{I}_x : A \rightarrow B\}_{x \in \mathbb{X}}$  of completely positive (CP) linear maps from  $\mathcal{B}(\mathcal{H}_A)$  to  $\mathcal{B}(\mathcal{H}_B)$ , such that  $\sum_x \mathcal{I}_x$  is trace-preserving (TP)

**interpretation**: expected probability of outcome  $x$  is  $p(x) = \text{Tr}[\mathcal{I}_x(\varrho)]$ , and corresponding post-measurement state is  $\frac{1}{p(x)} \mathcal{I}_x(\varrho)$

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## first part: the problem with measurement sharpness

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### sharp POVMs: conventional definition

**definition (folklore):** a POVM  $\mathbf{P} = \{P_x\}_{x \in \mathbb{X}}$  is called **sharp** whenever all its elements are projectors, i.e.,  $P_x P_{x'} = \delta_{x,x'} P_x$  for all  $x, x' \in \mathbb{X}$

**intuition:** sharp POVMs are “sharp” because

- orthogonal projectors are “pointed”
- they can be measured in a **repeatable**, “clear-cut” way

**sharpness as a resource:** Paul Busch already in 2005 envisioned a “resource theory of sharpness” proposing a class of **sharpness measures**; most recent work is by Liu and Luo (2022), and by Mitra (2022)

**question:** how can sharpness be “processed”?

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# POVM processing

POVMs can be transformed using

- a **quantum preprocessing**, i.e., a CPTP linear map  $\mathcal{E}$  such that  $\{P_x\}_{x \in \mathbb{X}} \mapsto \{Q_x\}_{x \in \mathbb{X}}$  with  $Q_x = \mathcal{E}^\dagger(P_x)$
- a **classical postprocessing**, i.e., a conditional distribution  $\mu(y|x)$  such that  $\{P_x\}_{x \in \mathbb{X}} \mapsto \{Q_y\}_{y \in \mathbb{Y}}$  with  $Q_y = \sum_x \mu(y|x) P_x$
- a **convex mixture** with another fixed POVM  $\mathbf{T} = \{T_x\}_{x \in \mathbb{X}}$ , i.e.,  $\{P_x\}_{x \in \mathbb{X}} \mapsto \{\lambda P_x + (1 - \lambda) T_x\}_{x \in \mathbb{X}}$ , with  $\lambda \in [0, 1]$
- a composition of the above

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## the problem

Which processings are sharpness-non-increasing?

- quantum preprocessings: can turn non-sharp into sharp  $\rightsquigarrow$  **ILLEGAL**
- classical postprocessings: can turn non-sharp into sharp  $\rightsquigarrow$  **ILLEGAL**
- convex mixtures: **legal** if  $\mathbf{T}$  is “maximally dull”, **but we need to characterize maximally dull POVMs first**

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## new definition: sharp<sup>#</sup> POVMs

### Definition

A given POVM  $\mathbf{P} = \{P_x\}_{x \in \mathbb{X}}$  is sharp<sup>#</sup> whenever the set

$$\text{range } \mathbf{P} := \left\{ \mathbf{p} \in \mathbb{R}_+^{|\mathbb{X}|} : \exists \varrho \text{ state}, p_x = \text{Tr}[\varrho P_x], \forall x \right\}$$

coincides with the entire probability simplex (“sharp”!) on  $\mathbb{X}$ .

It is dull<sup>#</sup> whenever  $\text{range } \mathbf{P}$  is a singleton.

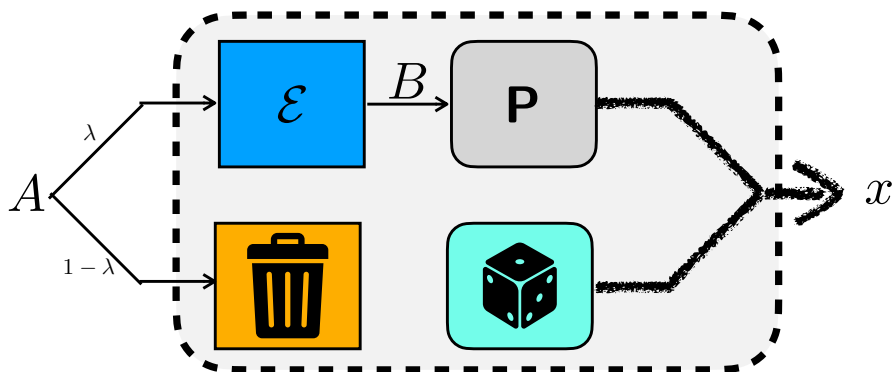
- sharp<sup>#</sup>  $\implies \dim \mathcal{H} \geq |\mathbb{X}| \implies$  nondegenerate observables are “canonical”
- sharp<sup>#</sup>  $\iff \forall x, \exists |\psi_x\rangle : P_x |\psi_x\rangle = |\psi_x\rangle$  ( $\implies P_x |\psi_{x'}\rangle = 0$  for  $x \neq x'$ )
- excluding null POVM elements, sharp  $\not\iff$  sharp<sup>#</sup>  $\iff$  repeatably measurable
- dull<sup>#</sup>  $\iff P_x \propto \mathbb{1}$ , for all  $x \in \mathbb{X}$

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## fuzzifying operations

for sharpness<sup>#</sup>:

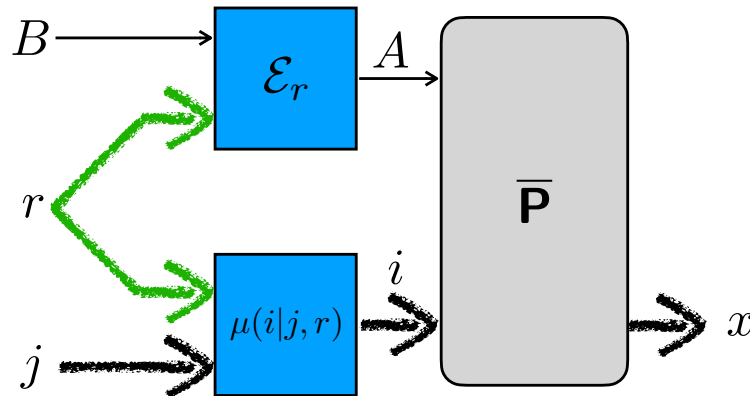
- quantum preprocessing: **LEGAL**
- convex mixture with any dull<sup>#</sup> POVM: **LEGAL**



$$P_x \longmapsto \lambda \mathcal{E}^\dagger(P_x) + (1 - \lambda) p(x) \mathbb{1}, \quad \forall x \in \mathbb{X}$$

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# fuzzifying operations as LOSR preprocessings of programmable POVMs



where

- $i, j \in \{0, 1, 2, \dots, |\mathbb{X}|\}$  label all possible programs
- $\bar{\mathbf{P}} = (\mathbf{P}, \mathbf{T}^{(1)}, \mathbf{T}^{(2)}, \dots, \mathbf{T}^{|\mathbb{X}|})$  is a programmable POVM with  $|\mathbb{X}| + 1$  program states, with  $\mathbf{T}^{(i)} = \{T_x^{(i)}\}_{x \in \mathbb{X}}$  denoting the deterministic POVMs, i.e.,  $T_x^{(i)} = \delta_{i,x} \mathbb{1}$

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## the sharpness<sup>#</sup> preorder

### Definition

given two POVMs  $\mathbf{P} = \{P_x\}_{x \in \mathbb{X}}$  and  $\mathbf{Q} = \{Q_x\}_{x \in \mathbb{X}}$ , we say that **P is sharper<sup>#</sup> than Q** ( $\mathbf{P} \succ^{\#} \mathbf{Q}$ ) whenever:

- there exists a fuzzifying operation transforming  $\mathbf{P}$  into  $\mathbf{Q}$
- equivalently: there exists a CPTP linear map  $\mathcal{E}$  such that

$$\mathbf{Q} \in \text{conv}\{\mathcal{E}^{\dagger}(\mathbf{P}), \mathbf{T}^{(1)}, \mathbf{T}^{(2)}, \dots, \mathbf{T}^{|\mathbb{X}|}\}$$

☞ note:  $\mathcal{E}$  trace-preserving  $\implies \mathcal{E}^{\dagger}$  unital  $\implies \mathcal{E}^{\dagger}(\mathbf{T}^{(i)}) = \mathbf{T}^{(i)}$

☞ this means that **all dull<sup>#</sup> are equivalent**: given any two dull<sup>#</sup> POVMs  $\mathbf{D}^{(1)}$  and  $\mathbf{D}^{(2)}$ , both  $\mathbf{D}^{(1)} \succ^{\#} \mathbf{D}^{(2)}$  and  $\mathbf{D}^{(2)} \succ^{\#} \mathbf{D}^{(1)}$  hold

☞ **question**: what about sharp<sup>#</sup> POVMs? are they also all equivalent?

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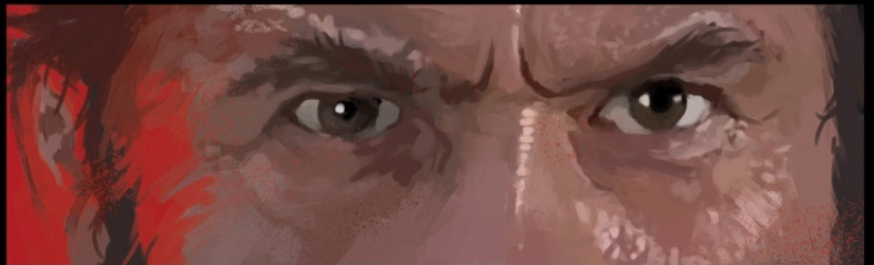
## the clean, the sharp<sup>#</sup>, and the dull<sup>#</sup>

clean



sharp

dull



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## the clean, the sharp<sup>#</sup>, and the dull<sup>#</sup>

### Definition

given two POVMs  $\mathbf{P} = \{P_x\}_{x \in \mathbb{X}}$  and  $\mathbf{Q} = \{Q_x\}_{x \in \mathbb{X}}$ , if a CPTP linear map  $\mathcal{E}$  exists such that  $Q_x = \mathcal{E}^\dagger(P_x) \forall x$ , we say that  $\mathbf{P}$  is **cleaner** than  $\mathbf{Q}$  ( $\mathbf{P} \succ \mathbf{Q}$ )

- POVMs  $\mathbf{P}$  such that, if  $\mathbf{Q} \succ \mathbf{P}$ , then also  $\mathbf{P} \succ \mathbf{Q}$ , are called **clean**
- **theorem**:  $\mathbf{P}$  clean  $\iff \mathbf{P} \succ \mathbf{Q}, \forall \mathbf{Q}$  (with same outcome set)
- **theorem**: clean  $\iff$  sharp<sup>#</sup>
- hence, **all sharp<sup>#</sup> are also equivalent**
- by the way: a POVM is dull<sup>#</sup> if and only if it is minimal for  $\succ$ , that is, if it satisfies  $\mathbf{P} \succ \mathbf{Q} \implies \mathbf{Q} \succ \mathbf{P}$

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## second part: the problem with instruments (in)compatibility

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### incompatibility

In quantum theory, some measurements necessarily exclude others.

This feature is what enables quantum algorithms, QKD protocols, violations of Bell's inequalities, etc.

Various formalizations:

- preparation uncertainty relations (e.g., Robertson)
- measurement uncertainty relations (e.g., Ozawa)
- incompatibility

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## compatible POVMs 1/2

### Definition

given a family  $\{\mathbf{P}^{(i)}\}_{i \in \mathbb{I}} \equiv \{P_x^{(i)}\}_{x \in \mathbb{X}, i \in \mathbb{I}}$  of POVMs, all defined on the same system  $A$ , we say that the family is *compatible*, whenever there exists

- a “mother” POVM  $\mathbf{O} = \{O_w\}_{w \in \mathbb{W}}$  on system  $A$
- a conditional probability distribution  $\mu(x|w, i)$

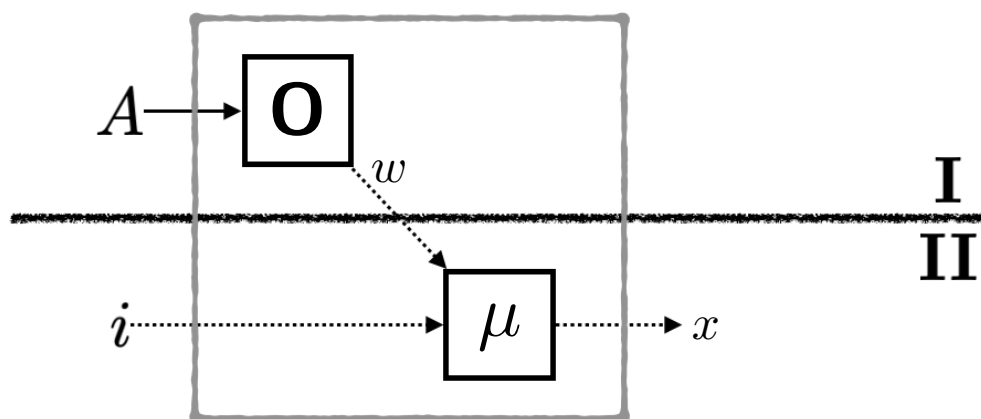
such that

$$P_x^{(i)} = \sum_w \mu(x|w, i) O_w ,$$

for all  $x \in \mathbb{X}$  and all  $i \in \mathbb{I}$ .

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## compatible POVMs 2/2



[F.B., E. Chitambar, W. Zhou; PRL 2020]

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## the first problem

While there is consensus on a single notion of compatibility for POVMs, the situation is less clear for instruments...

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## classical compatibility 1/2

### Definition (Heinosaari–Miyadera–Reitzner, 2014)

given a family of instruments  $\{\mathcal{I}_x^{(i)} : A \rightarrow B\}_{x \in \mathbb{X}, i \in \mathbb{I}}$ , we say that the family is *classically compatible*, whenever there exist

- a **mother instrument**  $\{\mathcal{H}_w : A \rightarrow B\}_{w \in \mathbb{W}}$
- a **conditional probability distribution**  $\mu(x|w, i)$

such that

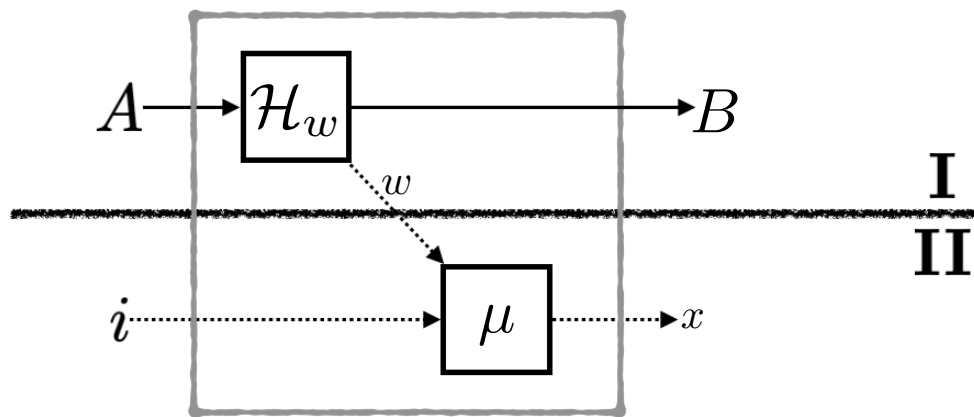
$$\mathcal{I}_x^{(i)} = \sum_w \mu(x|w, i) \mathcal{H}_w ,$$

for all  $x \in \mathbb{X}$  and all  $i \in \mathbb{I}$ .

we call this “classical” because it involves only **classical post-processings**, but it is also called “traditional” [Mitra and Farkas; PRA 2022].

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## classical compatibility 2/2



crucially:

- no shared entanglement and communication is classical
- communication goes only from **I** to **II**, i.e., the above is necessarily **II**→**I** non-signaling, see [Ji and Chitambar; PRA 2021]

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## marginalizing the mother

- without loss of generality (classical labels can be copied), compatible POVMs may be assumed to be recovered by marginalization, i.e.,

$$P_x^{(i)} = \sum_{x_j: j \neq i} O_{x_1, x_2, \dots, x_n}$$

- the notion of “parallel compatibility” for instruments lifts the above insight to the quantum outputs

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# parallel compatibility

## Definition (Heinosaari–Miyadera–Ziman, 2015)

given a family of instruments  $\{\mathcal{I}_x^{(i)} : A \rightarrow B_i\}_{x \in \mathbb{X}, i \in \mathbb{I}}$ , we say that the family is *parallelly compatible*, whenever there exist

- a **mother instrument**  $\{\mathcal{H}_w : A \rightarrow \otimes_{i \in \mathbb{I}} B_i\}_{w \in \mathbb{W}}$
- a **conditional probability distribution**  $\mu(x|w, i)$

such that

$$\mathcal{I}_x^{(i)} = \sum_w \mu(x|w, i) [\text{Tr}_{B_{i':i' \neq i}} \circ \mathcal{H}_w] ,$$

for all  $x \in \mathbb{X}$  and all  $i \in \mathbb{I}$ .

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# parallel compatibility VS classical compatibility

- parallel compatibility is able to go beyond no-signaling, hence, **parallel compatibility**  $\not\Rightarrow$  **classical compatibility**
- parallel compatibility has nothing to do with the “no information without disturbance” principle, because **non-disturbing instruments are never parallelly compatible**
- hence **classical compatibility**  $\not\Rightarrow$  **parallel compatibility**

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# bridging the two camps: q-compatibility

## Definition

given a family of instruments  $\{\mathcal{I}_x^{(i)} : A \rightarrow B_i\}_{x \in \mathbb{X}, i \in \mathbb{I}}$ , we say that the family is *q-compatible*, whenever there exist

- a **mother instrument**  $\{\mathcal{H}_w : A \rightarrow C\}_{w \in \mathbb{W}}$
- a **conditional probability distribution**  $\mu(x|w, i)$
- a **family of postprocessing channels**  $\{\mathcal{D}^{(x,w,i)} : C \rightarrow B_i\}_{x \in \mathbb{X}, w \in \mathbb{W}, i \in \mathbb{I}}$

such that

$$\mathcal{I}_x^{(i)} = \sum_w \mu(x|w, i) [\mathcal{D}^{(x,w,i)} \circ \mathcal{H}_w] ,$$

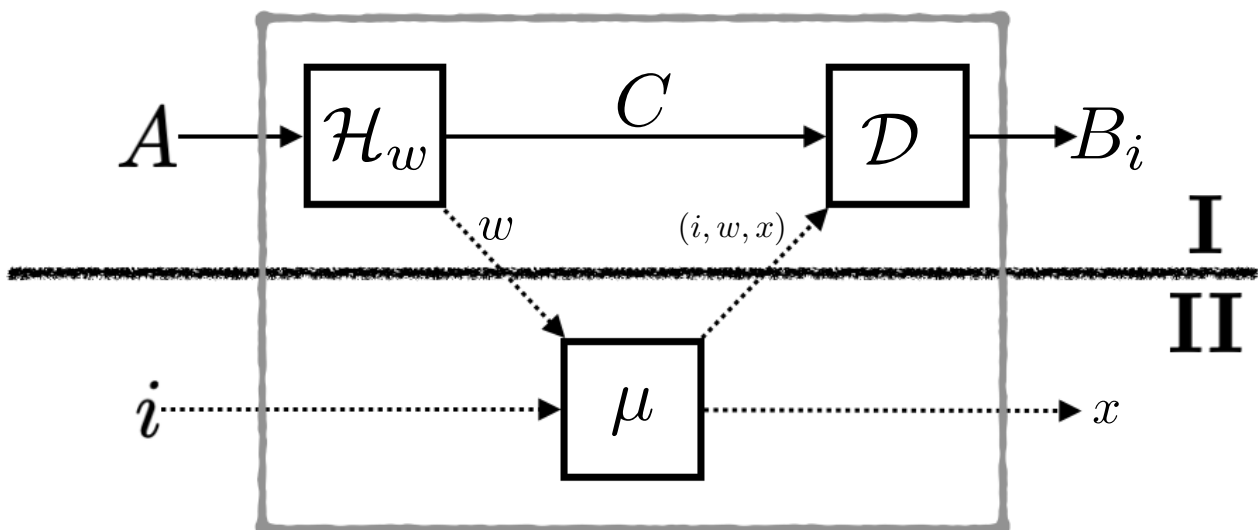
for all  $x \in \mathbb{X}$  and all  $i \in \mathbb{I}$ .

**classical compatibility:**  $C \equiv B_i$  and  $\mathcal{D}^{(x,w,i)} = \text{id}$

**parallel compatibility:**  $C \equiv \bigotimes_i B_i$  and  $\mathcal{D}^{(x,w,i)} = \text{Tr}_{B_{i'} : i' \neq i}$

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## q-compatibility as a circuit

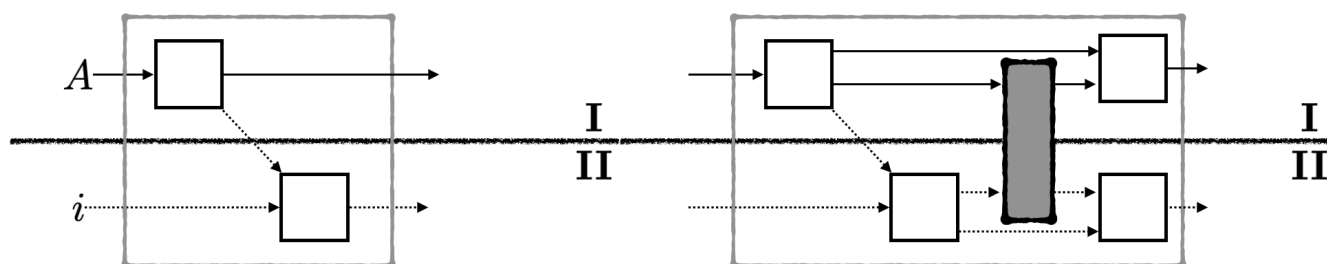


crucially:

- **no shared entanglement** and communication is **classical**
- only one interactive round **I**  $\rightarrow$  **II**  $\rightarrow$  **I**
- both classical and parallel compatibilities are special cases of q-compatibility

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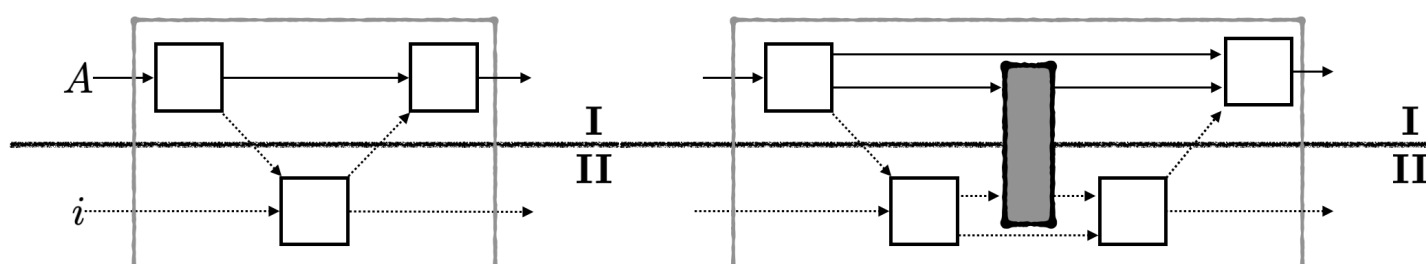
## free operations for classical incompatibility



- all classically compatible devices can be created for free
- if the initial device (the dark gray inner box) is classically compatible, the final device is also classically compatible

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## free operations for q-incompatibility



- all q-compatible devices can be created for free
- if the initial device (the dark gray inner box) is q-compatible, the final device is also q-compatible

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## the incompatibility preorder

given two families of instruments  $\{\mathcal{I}_x^{(i)} : A \rightarrow B_i\}_{x \in \mathbb{X}, i \in \mathbb{I}}$  and  $\{\mathcal{J}_y^{(j)} : C \rightarrow D_j\}_{y \in \mathbb{Y}, j \in \mathbb{J}}$ , we say

“ $\{\mathcal{I}_x^{(i)} : A \rightarrow B_i\}$  is more q-incompatible than  $\{\mathcal{J}_y^{(j)} : C \rightarrow D_j\}$ ”

whenever the former can be transformed into the latter by means of a free operation

$\rightsquigarrow$  this is now an instance of **statistical comparison**: a Blackwell-like theorem can be proved, and a complete family of monotones obtained

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## conclusion

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## take home messages

- **fuzzifying operations**: complete family of sharpness-non-increasing operations
- **sharpness is essentially a measure of classical communication capacity** (more precisely, signaling dimension)
- no need to argue about the “correct” definition of compatibility: **q-compatibility provides an overarching framework**
- **incompatibility is essentially quantum information transmission**, either in space (quantum channel) or in time (quantum memory)

thank you